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Some Results on Acyclic Chromatic Number of $\mathbf{M}^{\text {th }}$ \\ \section*{\title{
Some Results on Acyclic Chromatic Number of $\mathbf{M}^{\text {th }}$ Power of Paths
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## KEYWORDS

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#### Abstract

In this paper we discuss the graph power of paths, and some of their structural properties are derived. The acyclic chromatic number of $P_{n}^{m}$ and is derived both analytically and using their adjacency matrices. Also a relation between them is established.


## INTRODUCTION

Throughout this paper we are concerned only with finite, undirected simple graphs. Terms not defined here are used in the sense of Harary [3].

The notion of Graph power was introduced by Skiena [7] in 1990. The $\mathrm{k}^{\text {th }}$ power $\mathrm{G}^{\mathrm{k}}$ of an undirected graph G is another graph that has the same set of vertices, but in which two vertices are adjacent when their distance in $G$ is at most $k$ [2]. Powers of graphs are referred to using terminology similar to that for exponentiation of numbers: $\mathrm{G}^{2}$ is called the square of $G, G^{3}$ is called the cube of $G$, etc. Since a path of length two between vertices $u$ and $v$ exists for every vertex $w$ such that $\{u, w\}$ and $\{w, v\}$ are edges in $G$, the square of the adjacency matrix of $G$ counts the number of such paths. Similarly, the ( $u, v$ )th element of the $\mathrm{k}^{\text {th }}$ power of the adjacency matrix of $G$ gives the number of paths of length $k$ between vertices $u$ and $v$. It has been proved that $\operatorname{adj}\left(\mathrm{G}^{k}\right)=\sum_{i=1}^{k}[\operatorname{adj}(G)]^{i}$, where $\operatorname{adj}(\mathrm{G})$ is the adjacency matrix.

Graph colouring on the square of a graph may be used to assign frequencies to the participants of wireless communication networks so that no two participants interfere with each other at any of their common neighbours, and to find graph drawings with high angular resolution.

## PRELIMINARIES

The following basic definitions are taken from [3]. Graph colouring is an assignment of labels traditionally called "colours" to elements of a graph subject to certain constraints. The most common types of colourings are vertex colouring, edge colouring and face colouring. The vertex colouring is proper, if no two adjacent vertices are assigned the same colour. A proper vertex colouring of a graph is acyclic if every cycle uses at least three colours [4]. The acyclic chromatic number of $G$, denoted by $a(G)$, is the minimum colours required for its acyclic colouring.

The diameter $d$ of a graph is the maximum eccentricity of any vertex in the graph. That is, $d$ it is the greatest distance between any pair of vertices. It has been proved that, if a graph has diameter $d$ then its dth power is the complete graph [9].

## 1. ACYCLIC COLOURING OF mTH POWER OF AN nPATH $P_{n}$

The $m^{\text {th }}$ power of an $n$-path $P_{n}$ is a graph $P_{n}{ }^{m}$ with the same vertex set as $P_{n}$ in which two vertices are joined by
an edge if their distance in $P_{n}$ is atmost $m$.
1.1 Structural properties of $m^{\text {th }}$ power of the path $P_{n}$ for any $n$
The number of vertices in $P_{n}{ }^{m}$ is $p\left(P_{n}{ }^{m}\right)=n$
The number of edges in $P_{n}{ }^{m}$ is $q\left(P_{n}{ }^{m}\right)=m n-m(m+1)$, where $m, n \in N$ and $m<n-1$.

The minimum degree in $P_{n}{ }^{m}$ is $\delta\left(P_{n}{ }^{m}\right)=m, \quad n>1$ and $m<n$.
The maximum degree in $P_{n}{ }^{m}$ is $\Delta\left(P_{n}{ }^{m}\right)=n-1,\left\lfloor\frac{n}{2}\right\rfloor \leq m<n$.
For $m=2, P \frac{m}{n}$ is planar for all $n$.

### 1.2 Theorem

$\left.a\left(P_{n}^{n-1}\right)\right)=n$, for any $n \in N$.

## Proof

The eccentricity $\varepsilon(v)$ of the end vertices of $P_{n}$ are $n-1$. That is the diameter, $d=n-1$. So by the theorem 2.2 , for $d=n-1$, $P_{n}^{(n-1)}$ is isomorphic to $K_{n}$. Hence the theorem.

### 1.3 Theorem

$a\left(P_{n}^{m}\right)=m+1, \quad n \geq m+1$

## Proof

Case: 1 when $\mathrm{m}=1$
Then the theorem is trivially true
Case: 2 when $\mathrm{m}>1$
If $m=n-1$
by the theorem 2.3, $a\left(P_{n}^{n-1}\right)=n$, for any $n \in N$, the result is true.

If $m<n-1$
Let $V=\left\{v_{1} v_{2} v_{3}, \ldots v_{n}\right\}$ is the vertex set of the given graph $P_{n} m$. Consider the colour class $C=\left\{c_{1}, c_{2}, c_{3}, \ldots, c_{(m+1)}\right\}$. Since in $P_{n}{ }^{m}$ all the vertices with distance at most $m$ in $P_{n}$ are joined by an edge, we can find a complete sub graph $K_{m+1}$ with vertices $v_{1} v_{2} v_{3}, \ldots, v_{m+1}$, which are assigned the colours by the colour class $C$. Now the remaining $n-m-1$ vertices starting from $v_{(m+2)}, v_{(m+3)} v_{(m+4)}, \ldots$ are respectively assigned the colours $\mathrm{C}_{-1}, \mathrm{C}_{-2}, \mathrm{C}_{-3}, \ldots$ cyclically. Now the colouring is minimum as it contains $K_{m+1}$, minimum $m+1$ colours required for its
proper colouring. The colouring is acyclic, because for all $i \neq j$ the subgraph induced by the colour class <ci,cj> satisfies the relation $\varepsilon=\mathrm{v}-1$, which is the necessary and sufficient condition for a tree. Hence the theorem.


Figure 1: $a\left(P_{7}{ }^{3}\right)=4$

## 2. ADJACENCY MATRIX OF $P_{n}^{m}$

For $m<n$, the adiacencv matrix of $P_{n} m$ is an $n \times n$ matrix $A=[a i j]$, with $a i j=\left\{\begin{array}{l}1, \text { for }|i-j| \in(0, m] \\ 0, \text { else }\end{array}\right.$

### 2.1 Theorem

$a\left(P_{n}^{m}\right)=m i n T^{v d m}(v)+1, n \in N$, where $d^{m}(v)$ denote degree of vertex $v$ in the $m^{\text {th }}$ power of $P_{n}$.

## Proof:

For $m \leq n-1$, the graph $P_{n} m$ is obtained by joining each vertex of $P_{n}$ to nearby vertices which are at distance of at most $m$, so that the end vertices have degree $m$, which is the minimum degree. So by the theorem 1.3, $a\left(P_{n}{ }^{m}\right)=\delta+1$. Hence in the adjacency matrix of $P_{n}{ }^{m}, a\left(P_{n}{ }^{m}\right)=\min _{v} d^{m}(v)+1$, $n \in N$.

## Example

The adjacency matrix of $P_{12}{ }^{5}$ is the $12 \times 12$ matrix $A=[a i j]$, where $a_{i j}=\{(1$, for $|i-j| \in(0,5] 0$, else

$$
A=\left[\begin{array}{llllllllllll}
0 & 1 & 1 & 1 & 1 & 1 & 0 & 0 & 0 & 0 & 0 & 0 \\
1 & 0 & 1 & 1 & 1 & 1 & 1 & 0 & 0 & 0 & 0 & 0 \\
1 & 1 & 0 & 1 & 1 & 1 & 1 & 1 & 0 & 0 & 0 & 0 \\
1 & 1 & 1 & 0 & 1 & 1 & 1 & 1 & 1 & 0 & 0 & 0 \\
1 & 1 & 1 & 1 & 0 & 1 & 1 & 1 & 1 & 1 & 0 & 0 \\
1 & 1 & 1 & 1 & 1 & 0 & 1 & 1 & 1 & 1 & 1 & 0 \\
0 & 1 & 1 & 1 & 1 & 1 & 0 & 1 & 1 & 1 & 1 & 1 \\
0 & 0 & 1 & 1 & 1 & 1 & 1 & 0 & 1 & 1 & 1 & 1 \\
0 & 0 & 0 & 1 & 1 & 1 & 1 & 1 & 0 & 1 & 1 & 1 \\
0 & 0 & 0 & 0 & 1 & 1 & 1 & 1 & 1 & 0 & 1 & 1 \\
0 & 0 & 0 & 0 & 0 & 1 & 1 & 1 & 1 & 1 & 0 & 1 \\
0 & 0 & 0 & 0 & 0 & 0 & 1 & 1 & 1 & 1 & 1 & 0
\end{array}\right]
$$

Now using the theorem 6.1, $a\left(P_{n}^{m}\right)=\min _{v} d^{m}(v)+1$
Therefore, we have $a\left(P_{12}{ }^{5}\right)=\min _{v} d^{5}(v)+1=5+1=6$.
$a\left(P_{n}^{m}\right)=($ Number of integers in $(0, m])+1$.

## Example

In the previous example, $a\left(P_{12}{ }^{5}\right)=($ Number of integers in $(0,5])+1=5+1=6$.

## 3. RELATION BETWEEN ACYCLIC CHROMATIC NUMBERS OF C_n^m AND P_n^m

3.1 Observations
$a\left(C_{n}{ }^{m}\right)=2 a\left(P_{n}^{m}\right)-2$ when $m=\left\lfloor\frac{n}{2}\right\rfloor$ and for even $n \geq 3$
$a\left(C_{n}{ }^{m}\right)=2 a\left(P_{n}^{m}\right)-1$ when $m<\left\lfloor\frac{n}{2}\right\rfloor$ and for all $n \geq 3$

## CONCLUSIONS

In this paper, we propose some methods to determine the acyclic chromatic number of mth power of an $n$-path. We defined the acyclic chromatic number of this graphs in terms of adjacency matrices. This makes acyclic colouring of such graphs easy since the adjacency matrices of different graphs is a widely studied area in graph theory. A relatinn between the acyclic chromatic number of $\mathrm{C}_{n}{ }^{m}$ and $P\left\lfloor\frac{n}{2}\right\rfloor$ is also obtained.

### 2.2 Corollary

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