

Design and implementation of convolution encoder and viterbi decoder



Engineering

KEYWORDS : Convolutional encoder, Viterbi decoder, Verilog HDL, Viterbi Algorithm.

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ABSTRACT

Convolutional encoding is considered one of the forward error correction schemes. It is an essential component of wireless communication systems and mobile systems, which utilize some formulation of Convolutional encoding usually decoded via Viterbi decoders. A Viterbi decoder uses the Viterbi algorithm for decoding a bit stream that has been encoded using Forward error correction based on a Convolutional code. The maximum likelihood detection of a digital stream is possible by Viterbi algorithm. In this paper, we present a Convolutional encoder and Viterbi decoder with a constraint length of 9 and code rate of 1/2. This is realized using Verilog HDL. It is simulated and synthesized using Xilinx 13.1i.

Introduction:

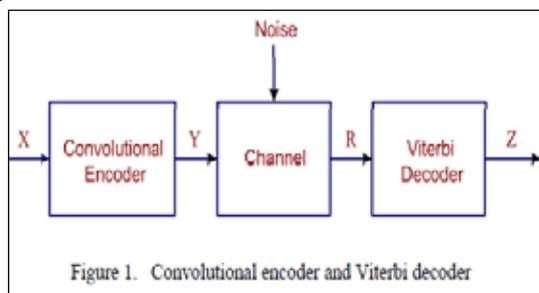
Convolution codes are used in a variety of systems including today's popular wireless standards (such as 802.11) and in satellite communications. standard for CDMA (Code Division Multiple Access), employs convolutional coding.

The Viterbi decoding algorithm was proposed and analyzed by Viterbi in 1967 [1]. It is widely used as a decoding technique for convolutional codes as well as the bit detection method in storage devices. Viterbi decoders currently find their use in more than one billion cell phones. The algorithm works by forming trellis structure, which is eventually traced back for decoding the received information. Convolutional encoding with Viterbi decoding is a powerful method for forward error correction. The Viterbi algorithm essentially

performs maximum likelihood decoding. However, it reduces the computational complexity by using trellis structure. Figure 1 shows the convolutional encoder and Viterbi decoder, which is used in the digital communication system. Here, X is the input data stream, which is given into the convolutional encoder and it produces the encoded data stream (Y). The encoded data stream (Y) is given to the channel in the presence of noise. Hence, it produces the noise added encoded data stream (R). Finally, data stream (R) is given to the Viterbi decoder that produces the estimated data stream (Z) applied at the input.

II. Convolution encoder

A convolutional encoder is a linear system. A binary convolutional encoder can be represented as a shift register. The outputs of the encoder are modulo 2 sums of the values in the certain register's cells. The input to the encoder is either the unencoded sequence (for non-recursive codes) or the unencoded sequence added with the values of some register's cells (for recursive codes). Convolutional codes can be systematic and non-systematic..



Convolutional codes are usually described using two parameters: the code rate and the constraint length. The code rate is expressed as a ratio of number of input symbols (k) into the channel encoder to the number of output symbols (n) by the channel encoder in a given cycle. Then, the code rate is expressed as, $r = k/n$ bits/symbol

al encoder. Convolutional encoder increases the length of the message sequence by adding redundant bits in order to increase the likelihood of detecting the transmitted sequence even if errors have occurred during transmission. Figure 2 shows the convolutional encoder of constraint length (K) = 9 and code rate (r) = 1/2. A convolutional encoder generates redundant bits by using modulo-2 convolutions. Hence, it is called as Convolutional encoder. If n modulo-2 adders are used, then it produces n outputs for each bit. The motivation of this paper is to realize a Viterbi decoder having Constraint length 9 and code rate 1/2 by using Xilinx 13.1i tools.

the state information of a convolutional encoder. In the state diagram, the state information of the convolutional encoder is shown in circles. Then, it is stored in the shift registers. Each new input information bit causes a transition from one state to another. The path information between the states has been denoted as x/c. Here, x represents the input information bits and c represents the output encoded bits. It is customary to begin convolutional encoding from the all zero state. For example, the input information sequence $x = \{1011\}$ (begin from the all zero state) leads to the state transition sequence $s = \{10, 01, 10, 11\}$ and produces the output encoded sequence $c = \{11, 10, 00, 01\}$.

III. Viterbi algorithm

Viterbi decoder is not a function of the number of symbols in the codeword sequence. The algorithm involves calculating a measure of similarity, or distance, between the received signal, at time t_i , and all the trellis paths entering each state at time t_i . The Viterbi algorithm removes from consideration those trellis paths that could not possibly be candidates for the maximum likelihood choice. When two paths enter the same state, the one having the best metric is chosen; this path is called the surviving path. This selection of surviving paths is performed for all the states. The decoder continues in this way to advance deeper into the trellis, making decisions by eliminating the least likely paths.

- In other words, it implies that the car first enter the stop state and then enter the

Reverse state. Hence, when we receive the information through the processes of forward, reverse and stop, we can safely interpret it as – forward, stop and reverse as this is a “maximum likelihood sequence”. The Viterbi algorithm uses the trellis diagram to compute the path metric value (accumulated distance) from the received sequence to the possible transmitted sequences. The total number of such trellis paths increases exponentially with the number of stages in the trellis. It causes potential complexity and memory problems. The Viterbi decoding algorithm has been classified into hard decision decoding and soft decision decoding. If the received signal is converted into two levels, either zero or one, it is called hard decision. If the input signal is quantized and processed for more than two levels, it is called soft decision. The soft decision decoding is expensive and requires large amount of memory than hard decision decoding. Hence, this work focuses on the hard decision decoding. Figure 4 shows the trellis diagram

for hard decision Viterbi decoding. When a sequence of data is received from the channel, it is desirable to estimate the original sequence that has been sent. The process of identifying such a sequence can be done using a diagram called 'trellis'.

The detection of the original stream can be described as finding the most probable path through the trellis. In the trellis diagram each node specifies an individual state at a given time and indicates a possible pattern of recently received data bits. The transition to a new state at the next timing cycle is indicated by each branch.

IV. VITERBI DECODER

Viterbi algorithm is used in the Viterbi decoder for decoding a bit stream that has been encoded using FEC based on a Convolutional code. Figure 5 shows the block diagram of Viterbi decoder [3]. It consists of the following functional units, namely, Branch Metric Unit, Path Metric Unit, Survivor Memory unit.

A. Branch Metric Unit

The comparison between received code symbol and expected code symbol is done by branch metric unit. It also counts the number of differing bits. It is the smallest unit in the Viterbi decoder. The measured value of the BMU can be the input decoding [4]. A branch metric unit's function is to calculate branch metrics, which are normed distances between every possible symbol in the code alphabet, and the received symbol.

Path Metric Unit

The partial path metric of each branch is computed by the use of two adders. The partial path metric is compared by the comparator and an appropriate branch is selected by the selector. The selector selects the smaller value and stored that value as the new path metric for each state. The corresponding bit decision is transferred to the SMU [4]. A path metric unit summarizes branch metrics to get metrics for $2K - 1$ paths, one of which can eventually be chosen as optimal. Every clock has some sequence it makes $2K - 1$ decisions, throwing off wittingly nonoptimal paths. The results of these decisions are written to the memory of a traceback unit.

Survivor Memory Unit

The Survivor Memory Unit receives the bit decision from the PMU. This will produce the decoded sequence. After finishing the SMU, it is important to perform the trace back module. When the trellis diagram is finished, the trace-back module will search the maximum likelihood path from the final state which is state0 to the beginning state which is state0.

V. IMPLEMENTATION

The Convolutional encoder and Viterbi decoder are implemented using Verilog HDL and the code has been developed under full-custom design. This implementation is complicated when using Verilog HDL compared to VHDL [2]. The Convolutional encoder encodes the given input sequence by modulo-2 adders. The operation is shown in fig. 2. The Viterbi decoder decodes the original input sequence by using Viterbi algorithm

From the trellis structure, we calculate the branch metric value and load that value. These values are added and compared. Then, the smallest value found in the process is the new path metric value. If all the states and the trellis stages are finished, then the trace back module is performed. It produces the decoded data (original input data).

VI. RESULTS AND DISCUSSION

The Convolutional encoder for the constraint length of $K=9$ and code rate of $r=1/2$ has been developed and the synthesis is carried

out. Table I shows the device utilization summary of the convolutional encoder. It has been simulated and the simulation result is shown in fig.

Typical input and output are as indicated below.
 Input bits X = 10010110011010010
 Encoded bits Y = 110110010110
 Encoded noise output R = 1101011001101010

fig.:Simulation results of Convolutional encoder

The Viterbi decoder has been developed using and the synthesis is carried out. Table II shows the device utilization summary of the Viterbi decoder. It has been simulated and the simulation result is shown in fig.

Typical input and output are as indicated below. Encoded noise
 Input bits X = 10010110011010010
 Encoded bits Y = 110110010110
 Encoded noise output R = 1101011001101010



fig.:Simulation results of Convolutional decoder

The Viterbi decoder has been developed using and the synthesis is carried out. Table II shows the device utilization summary of the Viterbi decoder. It has been simulated and the simulation result is shown in fig.

Typical input and output are as indicated below. Encoded noise
 input bits R = 1101011001101010
 Decoded output bits Z = 1001011001101001

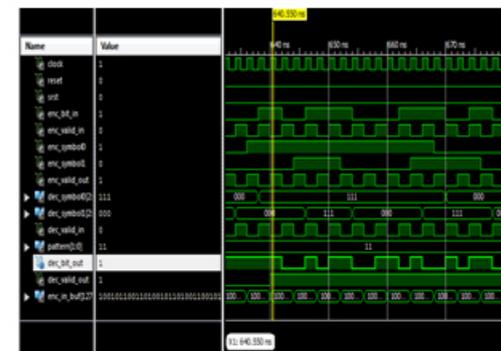


Fig: simulation results for decoder

VII. CONCLUSION

In this paper, we have presented the design and implementation of the Convolutional encoder and Viterbi decoder. This design has been simulated and synthesized using XILINX-ISE 13.1i for the constraint length of $K=9$ and code rate of $1/2$ input sequence. The given input sequence has been encoded by using convolution encoder and it is transmitted through the channel. Finally, the transmitted sequence is decoded by the Viterbi decoder and the estimated original sequence is produced.

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