

# A Novel Non-Monotone Self-Adaptive Trust Region Method for Unconstrained Optimization



## Mathematics

**KEYWORDS:** non-monotone strategy; self-adaptive trust region method; unconstrained optimization; global convergence.

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### ABSTRACT

*In this paper, we propose a new non-monotone self-adaptive trust region method for solving unconstrained optimization problem. Different from the usual trust region methods, our new algorithm uses a new rule to update the trust region radius. Theoretical analysis indicates that the new method preserves the global convergence under some reasonable assumptions.*

### 1. Introduction

Consider the following large unconstrained optimization problem:

$$\min f(x), \quad x \in R^n \tag{1.1}$$

where  $f(x) : R^n \rightarrow R$  is a twice continuously differentiable function.

For a given iteration point  $x_k$ , the trust region methods compute a trial step  $d_k$  by solving the following quadratic sub-problem:

$$\begin{aligned} \min \quad & q_k(d) = f_k + g_k^T d + \frac{1}{2} d^T B_k d, \\ \text{s.t.} \quad & \|d\| \leq \Delta_k, \end{aligned} \tag{1.2}$$

Where  $f_k = f(x_k)$ ,  $g_k = \nabla f(x_k)$ ,  $B_k \in R^{n \times n}$  is a symmetric matrix which is the Hessian matrix or its approximation of  $f(x)$  at the current point  $x_k$ ,  $\Delta_k > 0$  is the trust radius and  $\|\cdot\|$  denotes to the Euclidean norm. The ratio  $r_k$  between the actual reduction  $f(x_k) - f(x_{k+1})$  and the predicted reduction  $q_k(0) - q_k(d_k)$  plays a key role to decide whether  $d_k$  is acceptable or not and how to adjust the trust region radius. The iteration is said to be successful whenever  $r_k$  greater than a positive constant  $\eta_1$ . This leads us to the new point  $x_{k+1}$ , and the trust region radius is updated. If not, the iteration is unsuccessful, and the trial point is rejected. Generally, trust region radius update rule can be described as follow:

$$\Delta_{k+1} = \begin{cases} [\gamma_1 \Delta_k, \gamma_2 \Delta_k], & \text{if } r_k < \eta_1; \\ [\gamma_2 \Delta_k, \Delta_k], & \text{if } \eta_1 \leq r_k \leq \eta_2; \\ [\Delta_k, \infty), & \text{if } r_k \geq \eta_2. \end{cases}$$

(1.3)

Where the constants  $\gamma_1$ ,  $\gamma_2$ ,  $\eta_1$  and  $\eta_2$  satisfy

$$0 \leq \eta_1 < \eta_2 < 1, \quad 0 < \gamma_1 \leq \gamma_2 < 1 \tag{1.4}$$

In the trust region method, the difficulty is how to adjust the trust region radius. In order to choose an adaptive trust region radius at each iteration, many adaptive trust region methods have

been Studied in

In 1997, Sartenaer [3] introduced a strategy that can automatically determine the initial trust region radius. The strategy requires additional evaluations of the objective function. Zhanget al. [4] presented another strategy of updating the trust region radius. Their basic idea is originated from the following sub - problem in [5].

$$\begin{aligned} \min \quad & q_k(d) = g_k^T d + \frac{1}{2} d^T B_k d, \\ \text{s.t.} \quad & -\Delta_k \leq d_i \leq \Delta_k, \quad i = 1, 2, \dots, n, \end{aligned} \tag{1.5}$$

where  $\Delta_k = c^p (\|g_k\|/\gamma)$ ,  $0 < c < 1$ ,  $\gamma = \min(\|B_k\|, 1)$  and  $p$  is a nonnegative integer.

Hei [10] and Walmget al.[11] proposed self-adaptive update method, respectively. The trust region radius are product of so-called  $R$  -function and the  $\Lambda$  -function, that is,  $\Delta_{k+1} = R(r_k)\Delta_k$

and  $\Delta_{k+1} = \Lambda(r_k)\Delta_k$ . Recently, the  $L$  -function is introduced by Lu et al. in [12]. They presented

a new self-adaptive trust region method, in which the radius is  $\Delta_{k+1} = L(r_k)\Delta_k$ .

Grippio et al. [13] firstly proposed a non-monotone line search for Newton’s method. This algorithm accepts the step-size  $\alpha_k$  whether

$$f(x_k + \alpha_k d_k) \leq f(x_{l(k)}) + \beta \alpha_k \nabla f(x_k)^T d, \tag{1.6}$$

where  $\beta \in (0, \frac{1}{2})$ ,  $f(x_{l(k)}) = \max_{0 \leq j \leq m_k} f(x_{k-j})$ ,  $m_0 = 0$ ,  $0 \leq m_k \leq \min\{m_{k-1} + 1, M\}$  ( $k \geq 1$ ), and

$M \geq 0$  is an integer. It has been proved that the sequence  $\{f(x_k)\}$  is not increasing. Since then, many researchers [4-5] have exploited the non-monotone technique and a lot of numerical tests have showed that the non-monotone technique proposed by Grippo et al. [13] is efficient at some extent. In 1993, Deng et al. in [14] made some changes and applied it to the trust region method, and proposed a non-monotone trust region method for unconstrained optimization. Theoretical analysis and numerical results show that algorithms with non-monotone strategy are more effective than algorithms without it.

Zhang and Hager [15] proposed another non-monotone line search method, they replaced the maximum function value with an average of function values. In detail, their method finds a step-size  $\alpha_k$  satisfying the following condition:

$$f(x_k + \alpha_k d_k) \leq C_k + \beta \alpha_k \nabla f(x_k)^T d, \tag{1.7}$$

where

$$C_k = \begin{cases} f(x_k), & k = 0, \\ \frac{\eta_{k-1} Q_{k-1} C_{k-1} + f(x_k)}{Q_k}, & k \geq 1, \end{cases} \quad Q_k = \begin{cases} 1, & k = 0, \\ \eta_{k-1} Q_{k-1} + 1, & k \geq 1, \end{cases} \tag{1.8}$$

and  $\eta_{k-1} \in [\eta_{\min}, \eta_{\max}]$ ,  $\eta_{\min} \in [0, 1)$  and  $\eta_{\max} \in [\eta_{\min}, 1)$  are two chosen parameters. Numerical

results showed that this non-monotone technique was superior to (1.6). Then, this non-monotone was applied to the trust region methods[16,17].In addition, many non-monotone adaptive trust region methods have been proposed in [6-9].

Inspired by the ideas introduced above, we use the  $L$ -function to update the radius, then applied it to the trust region method with non-monotone strategy proposed by Zhang and Hager [15].The purpose of this paper is to present a new non-monotone adaptive trust region method.

The rest of the paper is organized as follows. In Section 2, we describe  $L$ -function and our new non-monotone self-adaptive trust region algorithm. The global convergence properties of this novel algorithm are given in Section 3. Finally, some conclusions are summarized in Section 4.

## 2. $L$ -function and Algorithm

The  $L$ -function rule proposed by Lu [14] can be described as follows:

$$\Delta_{k+1} = L(r_k)\Delta_k \tag{2.1}$$

where the  $L$ -function  $L(r_k)$  is chosen as

$$L(r_k) = \begin{cases} c_1 + (c_2 - c_1) \exp(r_k), & \text{if } r_k \leq 0, \\ \frac{1 - \beta_1 \exp(\eta_2)}{1 - \exp(\eta_2)} - \frac{(1 - \beta_1) \exp(\eta_2)}{1 - \exp(\eta_2)} \exp((r_k - \eta_2)), & \text{if } 0 < r_k < \eta_2, \\ \beta_2, & \text{if } \eta_2 \leq r_k < 2 - \eta_2, \\ \beta_3 + (\beta_2 - \beta_3) \exp(-(\frac{r_k + \eta_2 - 2}{\eta_2 - 2})^2), & \text{if } r_k > 2 - \eta_2, \end{cases} \tag{2.2}$$

where  $\beta_1, \beta_2, \beta_3, c_1, c_2$  and  $\eta_2$  are constants.

Now describe the new non-monotone self-adaptive trust region algorithm with the new radius update rule.

When we obtain  $d_k$ , then the ratio  $r_k$  is computed by

$$r_k = \frac{Ared_k}{Pred_k} = \frac{C_k - f(x_k + d_k)}{q_k(0) - q_k(d_k)}, \tag{2.3}$$

### Algorithm 2.1

Step 1. Given  $x_0 \in R^n, \Delta_0 > 0, B_0 \in R^{n \times n}, 0 < \eta_1 < \eta_2 < 1, 0 < c_1 < c_2 < 1, \varepsilon \geq 0,$

$0 < \beta_1 \leq \beta_3 < 1 \leq \beta_2,$  set  $k := 0.$

Step 2. Compute  $g_k$ . If  $\|g_k\| \leq \varepsilon$ , stop. Otherwise, go to Step 3.

Step 3. Solve the sub-problem (1.2) for  $d_k$ . Compute  $C_k, Ared_k, Pred_k$  and  $r_k$ .

Step 4. If  $r_k > \eta_1$ , set  $x_{k+1} = x_k + d_k$ , otherwise,  $x_{k+1} = x_k$ .

Step 5. Update the trust region radius  $\Delta_{k+1}$  by (2.1) and (2.2).

Step 6. Compute  $g_{k+1}$  and  $B_{k+1}$ , and set  $k := k + 1$ , go to Step 2.

### 3. Convergence

In this section, we will prove the global convergence properties of Algorithm 2.1. The following assumptions are necessary to analyze the convergence properties.

(H1) The level set  $L(x_0) = \{x \in R^n \mid f(x) \leq f(x_0)\}$  is bounded for any given  $x_0 \in R^n$ .

(H2) The matrix  $B_k$  is uniformly bounded, i.e., there exists a constant  $M_0 > 0$ , such that, for all  $k$ ,

$$\|B_k\| \leq M_0.$$

**Lemma 3.1.** If  $d_k$  is the solution to sub-problem (1.2), then

$$Pred_k = q_k(0) - q_k(d_k) \geq \frac{1}{2} \|g_k\| \min \left\{ \Delta_k, \frac{\|g_k\|}{\|B_k\|} \right\}. \tag{3.1}$$

$$g_k^T d_k \leq -\frac{1}{2} \|g_k\| \min \left\{ \Delta_k, \frac{\|g_k\|}{\|B_k\|} \right\} \tag{3.2}$$

Proof. From Lemma 13.3.1 in [18], we know (3.1) holds. And from (3.1), we can see

$$Pred_k = -g_k^T d_k - \frac{1}{2} d_k^T B_k d_k \geq \frac{1}{2} \|g_k\| \min \left\{ \Delta_k, \frac{\|g_k\|}{\|B_k\|} \right\}$$

Consider the above inequality and the fact  $d_k^T B_k d_k > 0$ , (3.2) obviously holds. Therefore, the lemma is true.

**Lemma 3.2.** Let  $\{x_k\}$  be the sequence generated by Algorithm 2.1. For any fixed  $k \geq 0$ , we have

$$f_{k+1} \leq C_{k+1} \leq C_k \tag{3.3}$$

Proof. Let  $k \geq 0$  be an arbitrary fixed integer. By the definition of  $r_k$  and  $r_k > \eta_1$ , we have

$$\begin{aligned} C_k - f_{k+1} &\geq \eta_1 Pred_k = \eta_1 (q_k(0) - q_k(d_k)) \\ &\geq \frac{\eta_1}{2} \|g_k\| \min \left\{ \Delta_k, \frac{\|g_k\|}{\|B_k\|} \right\} \geq 0 \end{aligned} \tag{3.4}$$

Thus,  $C_k \geq f_{k+1}$ . Then, by the definition of  $C_k$ , we obtain that

$$C_k = \frac{\eta_{k-1} Q_{k-1} C_{k-1} + f_k}{Q_k} \geq \frac{\eta_{k-1} Q_{k-1} f_k + f_k}{Q_k} = f_k.$$

So

$$C_k \geq f_k. \tag{3.5}$$

On the other hand, we have

$$C_{k+1} = \frac{\eta_k Q_k C_k + f_{k+1}}{Q_{k+1}} \leq \frac{\eta_k Q_k C_k + C_k}{Q_{k+1}} = C_k. \tag{3.6}$$

From (3.5) and (3.6), Lemma 3.2 holds.

**Lemma 3.3.** Suppose that the sequence  $\{x_k\}$  is generated by Algorithm 2.1. The algorithm is well defined.

Proof .The process is similar to Lemma 3.4 in [19].

**Lemma 3.4.**Suppose that(H1)-(H2) hold. Then

$$|f(x_k + d_k) - q(d_k)| = O(\|d_k\|^2) \tag{3.7}$$

Proof. From the Taylor expansion, the lemma is true.

**Lemma 3.5.**(See Lemma 3.6 in [19]) Suppose that Assumption 1 holds, and there is a positive number  $\tau > 0$  such that  $\|g_k\| \geq \tau$  for all  $k$ , then there exists a  $\bar{\Delta} > 0$ , such that for all  $k$ , we have  $\Delta_k \geq \bar{\Delta}$ .

**Theorem 3.6.**Suppose that (H1)-(H2) hold. Let the sequence  $\{x_k\}$  generated by Algorithm 2.1, then we have

$$\liminf_{k \rightarrow \infty} \|g_k\| = 0. \tag{3.8}$$

Proof.We assume that Formula (3.8) is not true,that is, there exists a positive constant  $\tau > 0$ , such that

$$\|g_k\| \geq \tau \text{ for all } k. \tag{3.9}$$

By Lemma 3.1, Lemma 3.2, we have

$$f_{k+1} \leq C_k - \eta_1 \text{Pred}_k \leq C_k - \frac{1}{2} \eta_1 \|g_k\| \min\{\Delta_k, \frac{\|g_k\|}{\|B_k\|}\} \tag{3.10}$$

Then

$$\begin{aligned} C_{k+1} &= \frac{\eta_k Q_k C_k + f_{k+1}}{Q_{k+1}} \leq \frac{\eta_k Q_k C_k + C_k - \frac{1}{2} \eta_1 \|g_k\| \min\{\Delta_k, \frac{\|g_k\|}{\|B_k\|}\}}{Q_{k+1}} \\ &= C_k - \frac{\frac{1}{2} \eta_1 \|g_k\| \min\{\Delta_k, \frac{\|g_k\|}{\|B_k\|}\}}{Q_{k+1}} \\ &\leq C_k - \frac{\frac{1}{2} \eta_1 \tau \min\{\bar{\Delta}, \frac{\tau}{M_0}\}}{Q_{k+1}} \end{aligned} \tag{3.11}$$

Which implies

$$C_k - C_{k+1} \geq \frac{\frac{1}{2} \eta_1 \tau \min\{\bar{\Delta}, \frac{\tau}{M_0}\}}{Q_{k+1}} \tag{3.12}$$

It follows that lemma 3.2 that  $f_k \leq C_k$  for all  $k$  and  $\{C_k\}$  is decreasing. Thus, by assumptions, lemma 3.3 and Formula (3.21), we have that  $\{f_k\}$  is bounded below. Therefore,  $\{C_k\}$  is convergent. From (3.12) that

$$\sum_{k=0}^{\infty} \frac{\min\{\bar{\Delta}, \frac{\tau}{M_0}\}}{Q_{k+1}} < \infty \tag{3.13}$$

By (1.8), we know  $Q_0 = 1$  and  $\eta_k \in [0, 1)$ , we have

$$Q_{k+1} = 1 + \sum_{i=0}^k \prod_{m=0}^i \eta_{k-m} \leq 1 + \sum_{i=0}^k \eta_{\max}^{i+1} \leq \sum_{i=0}^{\infty} \eta_{\max}^i = \frac{1}{1 - \eta_{\max}} \tag{3.14}$$

Set  $\min\{\bar{\Delta}, \frac{\tau}{M_1}\} = \lambda$ . Formula (3.13) can be written as

$$\sum_{k=0}^{\infty} k(1 - \eta_{\max})\lambda < \infty \tag{3.15}$$

we know  $\sum_{k=0}^{\infty} k(1 - \eta_{\max})\lambda$  is not convergent. This is a contradiction with Formula (3.15).

Theorem 3.6 has been proved.

### 4. Conclusions

In this paper, we present a new non-monotone self-adaptive trust region method. And the form of the new method is very simple. Under some mild conditions, we proved the global convergence result of the proposed method.

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